

## Enhanced Compressive Sensing and More

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CS	ECS	Rice L1	FPC	FTVd	Numerical Results
Outline					

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#### Collaborators

- Wotao Yin, Elaine Hale
- Students: Yilun Wang, Junfeng Yang
- Acknowledgment: NSF DMS-0442065

#### The Outline of the Talk

- Compressive Sensing (CS): when  $\ell_0 \Leftrightarrow \ell_1$ ?
- An accessible proof and Enhanced CS
- Rice L1-Related Optimization Project
- An L1 Algorithm: FPC
- A TV Algorithm: FTVd
- Numerical Results



## Compressive Sensing (CS)

#### Recover sparse signal from incomplete data:

- Unknown signal  $x^* \in \mathbb{R}^n$
- Measurements:  $b = Ax^* \in \mathbb{R}^m$ , m < n
- $x^*$  is sparse: #nonzeros  $||x^*||_0 < m$

#### 1 Solution to 2 Problems?

- $\ell_0$ -Prob: min{ $||x||_0 : Ax = b$ }  $\Rightarrow$  sparsest solution (hard)
- $\ell_1$ -Prob: min{ $||x||_1 : Ax = b$ }  $\Rightarrow$  lin. prog. solution (easy)
- **Recoverability:** When does the same *x*<sup>\*</sup> solve both?



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When does the following happen? ( $b = Ax^*$ )

 $\{x^*\} = \arg\min\{\|x\|_0 : Ax = b\} = \arg\min\{\|x\|_1 : Ax = b\}$ 

Answer: For a random  $A \in \mathbb{R}^{m \times n}$ ,

 $\|x^*\|_0 < \frac{c \cdot m}{\log(n/m)}.$ 

- Candes-Romberg-Tao, Donoho et al, 2005
- Rudelson-Vershynin, 2005, 2006
- Baraniuk-Davenport-DeVore-Wakin, 2007 .....





Theoretical guarantees available:

 $\min\{\|\Phi x\|_1 : Ax = b\}, \quad \min\{\|\Phi x\|_1 : Ax = b, x \ge 0\}$ (Donoho-Tanner 2005, Z 2005)

#### What about these convex models?

$$\min\{\|\Phi x\|_{1} + \mu TV(x) : Ax = b\}$$
$$\min\{\|\Phi x\|_{1} + \mu \|x - \hat{x}\| : Ax = b\}$$
$$\min\{\|\Phi x\|_{1} : Ax = b, Bx \le c, x \in [I, u]\}$$

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#### When is $\ell_0 \Leftrightarrow \ell_1$ ?

- Most analyses are based on the notion of RIP: —Restricted Isometry Property
- Or based on "counting faces" of polyhedrons
- Derivations are quite involved and not transparent
- Generalize CS analysis to more models?

#### A simpler, gentler, more general analysis?

Yes. Using Kashin-Garnaev-Gluskin (KGG) inequality.

(Extension to Z, CAAM Report TR05-09)



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#### $\ell_1$ -norm vs. $\ell_2$ -norm:

$$\sqrt{n} \geq \frac{\|\boldsymbol{v}\|_1}{\|\boldsymbol{v}\|_2} \geq 1, \quad \forall \boldsymbol{v} \in \mathbb{R}^n \setminus \{\mathbf{0}\}$$

However,  $\|v\|_1/\|v\|_2 \gg 1$  in most subspaces of  $\mathbb{R}^n$ .

#### Theorem: (Kashin 77, Garnaev-Gluskin 84)

Let  $A \in \mathbb{R}^{m \times n}$  be iid Gaussian. With probability  $> 1 - e^{-c_1(n-m)}$ ,

$$\frac{\|\boldsymbol{v}\|_1}{\|\boldsymbol{v}\|_2} \geq \frac{\boldsymbol{c}_2\sqrt{m}}{\sqrt{\log(n/m)}}, \quad \forall \boldsymbol{v} \in \textit{Null}(\textit{A}) \setminus \{0\}$$

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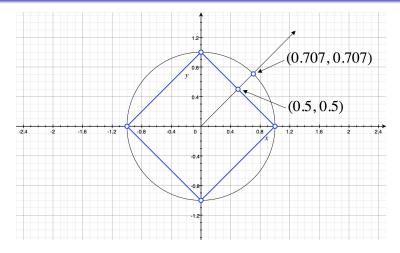
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where  $c_1$  and  $c_2$  are absolute constants.



## A Picture in 2D



In most subspaces,  $\|v\|_1 / \|v\|_2 \ge 0.8 * \sqrt{2} > 1.1$ 



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## CS ECS Rice L1 FPC FTVd Numerical Results Sparsest Point vs. $\ell_p$ -Minimizer, $p \in (0, 1]$

When does the following hold on  $C \subset \mathbb{R}^n$ ?

$$\{x^*\} = \arg\min_{x \in C} \|x\|_0 = \arg\min_{x \in C} \||x|^p\|_1$$

This means: (i) " $\ell_0 \Leftrightarrow \ell_p$ " on *C*, (ii) uniqueness of  $x^*$ .

A Sufficient Condition — entirely on sparsity

$$\sqrt{\|x^*\|_0} < rac{1}{2} rac{\||v|^{
ho}\|_1}{\||v|^{
ho}\|_2}, \ \ \forall \, v \in (\mathcal{C} - x^*) \setminus \{0\}$$

(10-line, elementary proof skipped)



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CS ECS Rice L1 FPC FTVd Numerical Results
Recoverability Proved and Generalized

For 
$$C = \{x : Ax = b, x \in S\}$$
,

$$C - x^* = Null(A) \cap (S - x^*), \quad \forall S \subset \mathbb{R}^n$$

$$[\ell_0 \Leftrightarrow \ell_p] \leftarrow \|x^*\|_0^{\frac{1}{2}} < \frac{1}{2} \frac{\||v|^p\|_1}{\||v|^p\|_2}, \forall v \in Null(A) \cap (S - x^*) \setminus \{0\}$$

#### For a Gaussian random A, by GKK

$$[\ell_0 \Leftrightarrow \ell_p]$$
 on  $C \stackrel{\text{h.p.}}{\longleftarrow} \|x^*\|_0 < \frac{c(p) \cdot m}{\log(n/m)}$ 

(Stability results also available for noisy data)



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ECS: with prior information  $x \in S$ 

$$\min\{\|x\|_1: Ax = b, x \in S\}$$

We have shown ECS recoverability is at least as good as CS.

More prior information (beside nonnegativity)?

$$\min\{\|\boldsymbol{x}\|_1 + \mu \operatorname{TV}(\boldsymbol{x}) : \boldsymbol{A}\boldsymbol{x} = \boldsymbol{b}\} \Rightarrow \boldsymbol{S} = \{\boldsymbol{x} : \operatorname{TV}(\boldsymbol{x}) \le \delta\}$$

$$\min\{\|x\|_1 + \mu\|x - \hat{x}\| : Ax = b\} \Rightarrow S = \{x : \|x - \hat{x}\| \le \delta\}$$

..... and many more possibilities.

More ECS models, more algorithmic challenges for optimizers.



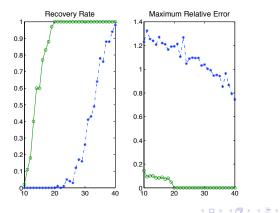
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Unknown signal  $x^*$  close to a prior sparse  $x_p$ :

ECS: 
$$\min\{||x||_1 : Ax = b, ||x - x_p||_1 \le \delta\}$$

With 10% differences in supports and nonzero values,





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Computational & Applied Math. Dept. in Engineering School:

- Y. Z., Wotao Yin (Elaine Hale, left)
- Students

#### Optimization Algorithmic Challenges in CS

- Large-scale, (near) real-time processing
- Dense matrices, non-smooth objectives
- Traditional (simplex, interior-point) methods have trouble.

#### Can convex optimization be practical in CS?



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Convex Optimization is generally more robust w.r.t noise.

### Is it too slow for large-scale applications?

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#### In many cases, it is faster than other approaches.

- Solution sparsity helps.
- Fast transforms help.
- Structured random matrices help.
- Efficient algorithms can be built on Av and  $A^Tv$ .
- Real-time algorithms are possible for problems with special structures (like MRI).
- 2 examples from our work: FPC and FTVd

## CS ECS Rice L1 FPC FTVd Numerical Results Forward-Backward Operator Splitting

Derivation (since 1950's):

$$\begin{aligned} \min \|x\|_{1} + \mu f(x) &\Leftrightarrow \quad \mathbf{0} \in \partial \|x\|_{1} + \mu \nabla f(x) \\ &\Leftrightarrow \quad -\tau \mu \nabla f(x) \in \tau \partial \|x\|_{1} \\ &\Leftrightarrow \quad x - \tau \mu \nabla f(x) \in x + \tau \partial \|x\|_{1} \\ &\Leftrightarrow \quad (I + \tau \partial \| \cdot \|_{1}) x \ni x - \tau \mu \nabla f(x) \\ &\Leftrightarrow \quad \{x\} \ni (I + \tau \partial \| \cdot \|_{1})^{-1} (x - \tau \mu \nabla f(x)) \\ &\Leftrightarrow \quad x = shrink(x - \tau \nabla f(x), \tau / \mu) \end{aligned}$$

#### Equivalence to Fixed Point

$$\min \|x\|_1 + \mu f(x) \iff x = Shrink(x - \tau \nabla f(x), \tau/\mu)$$



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$$\min_{x} \|x\|_1 + \mu f(x)$$

Algorithm:

$$\mathbf{x}^{k+1} = Shrink(\mathbf{x}^k - \tau 
abla f(\mathbf{x}^k), au/\mu)$$

where

$$Shrink(y, t) = y - Proj_{[-t,t]}(y)$$

- A "first-order" method follows from FB-operator splitting
- Discovered in signal processing by many since 2000's
- Convergence properties analyzed extensively



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## CS ECS Rice L1 FPC FTVd Numerical Results New Convergence Results (Hale, Yin & Z, 2007)

#### How can solution sparsity help a 1st-order method?

• Finite Convergence: for all but a finite # of iterations,

$$x_j^k = 0, \qquad ext{if } x_j^* = 0$$
  
 $\operatorname{sign}(x_j^k) = \operatorname{sign}(x_j^*), \qquad ext{if } x_j^* 
eq 0$ 

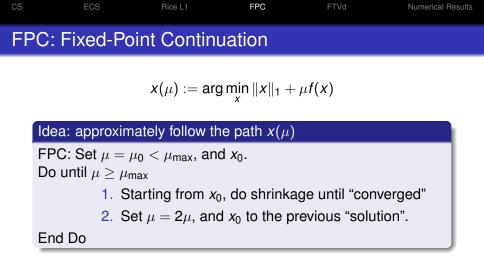
• q-linear rate depending on "reduced" Hessian:

$$\limsup_{k \to \infty} \frac{\|x^{k+1} - x^*\|}{\|x^k - x^*\|} \le \frac{\kappa(H_{EE}^*) - 1}{\kappa(H_{EE}^*) + 1}$$

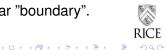
where  $H_{EE}^*$  is a sub-Hessian of *f* at  $x^*$  ( $\kappa(H_{EE}^*) \le \kappa(H^*)$ ), and  $E = \operatorname{supp}(x^*)$  (under a regularity condition).

The sparser  $x^*$  is, the faster the convergence.



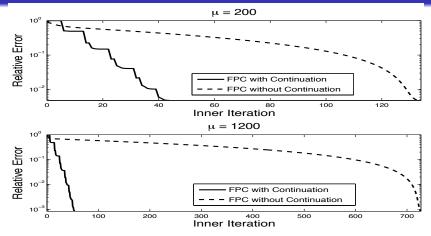


- Smaller  $\mu \rightarrow$  sparser  $x(\mu) \rightarrow$  faster convergence
- Converges is also fast for larger  $\mu$  due to 'warm starts".
- Generally effective, may slow down near "boundary".





### Continuation Makes It Kick



(Numerical comparison results in Hale, Yin & Z 2007)

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Fully random matrices are computationally costly.

Random Kronicker Products

$$A = A_1 \otimes A_2 \in \mathbb{R}^{m \times n},$$

where we only store  $A_1$  ( $m/p \times n/q$ ) and  $A_2$  ( $p \times q$ ).

$$Ax = \operatorname{vec}\left(A_2 X A_1^T\right)$$

also requires much less computation.

- Trial: n = 1M, m = 250k,  $||x^*||_0 = 25$ k, p = 500, q = 1000.
- 2 (500 by 1000) matrices stored (reduction of 0.5M times).
- FPC solved the problem in 47s on a PC.



# CS ECS Rice L1 FPC FTVd Numerical Results Total Variation Regularization

Discrete total variation (TV) for an image u:

 $TV(u) = \sum \|D_i u\|$  (sum over all pixels)

(1-norm of the vector of 2-norms of discrete gradients)

- Advantage: able to capture sharp edges
- Rudin-Osher-Fatemi 1992, Rudin-Osher 1994
- Recent Survey: Chan and Shen 2006

Non-smoothness and non-linearity cause computational difficulties. Can TV be competitive in speed with others (say, Tikhonov-like regularizers)?



## CS ECS Rice L1 FPC FTVd Numerical Results Fast TV deconvolution (FTVd)

(TV + L2) 
$$\min_{u} \sum \|D_{i}u\| + \frac{\mu}{2} \|Ku - f\|^{2}$$

Introducing  $w_i \in \mathbb{R}^2$  and a quadratic penalty (Courant 1943):

$$\min_{u,w} \sum \left( \|w_i\| + \frac{\beta}{2} \|w_i - D_i u\|^2 \right) + \frac{\mu}{2} \|\mathcal{K}u - f\|^2$$

In theory,  $u(\beta) \rightarrow u^*$  as  $\beta \rightarrow \infty$ . In practice,  $\beta = 200$  suffices.

#### Alternating Minimization:

- For fixed u,  $w_i$  can be solved by a 2D-shrinkage.
- For fixed w, quadratic can be minimized by 3 FFTs.

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(Wang, Yang, Yin &Z, 2007, 2008)
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## CS ECS Rice L1 FPC FTVd Numerical Results FTVd

FTVd is a long-missing member of the half-qudratic class (Geman-Yang 95), using a 2D Huber-like approximation.

#### FTVd Convergence

- Finite convergence for  $w_i^k \rightarrow 0$  (sparsity helps).
- Strong *q*-linear convergence rates for the others
- Rates depend on submatrices (sparsity helps).
- Continuation accelerates practical convergence.

#### FTVd Performance

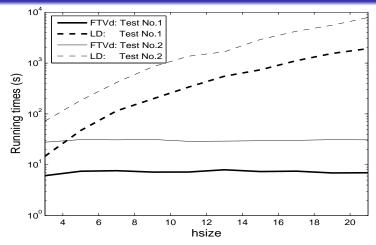
- Orders of magnitude faster than Lagged Diffusivity.
- Comparable speed with Matlab deblurring, with better quality. TV models has finally caught up in speed.



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## FTVd vs. Lagged Diffusivity



(Test 1: Lena 512 by 512; Test 2: Man 1024 by 1024)



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### FTVd vs. Others

Blurry&Noisy. SNR: 5.19dB



FTVd:  $\beta = 2^5$ , SNR: 12.58dB, t = 1.83s



deconvwnr: SNR: 11.51dB, t = 0.05s

deconvreg: SNR: 11.20dB, t = 0.34s

FTVd:  $\beta = 2^7$ , SNR: 13.11dB, t = 14.10s







## FTVd Extensions

Multi-channel Image Deblurring (paper forthcoming)

- cross-channel or within-channel blurring
- a "small" number of FFTs per iteration
- convergence results have been generalized
- TV+L<sup>1</sup> deblurring models (codes hard to find)

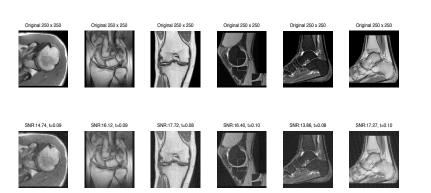
#### Other Extensions

- higher-order TV regularizers (reducing stair-casing)
- multi-term regularizations multiple splittngs
- locally weighted TV → weighted shrinkage
- reconstruction from partial Fourier coefficients (MRI)

 $\min_{u} TV(u) + \lambda \|\Phi u\|_1 + \mu \|\mathcal{F}_{\rho}(u) - f_{\rho}\|^2$ 

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## MRI Construction from 15% Coefficients



250 by 250 Images: time  $\leq$  0.1s on a PC (3 GHz Pentium D).



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## Color Image Deblurring

Original image: 512x512



Blurry & Noisy SNR: 5.1dB.



deconvlucy: SNR=6.5dB, t=8.9



deconvreg: SNR=10.8dB, t=4.4 deconvwnr: SNR=10.8dB, t=1.4



MxNopt: SNR=16.3dB, t=1.6



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Comparison to Matlab Toolbox:  $512 \times 512$  Lena



#### Take-Home Messages

- CS recoverability can be proved in 1 page via KGG.
- Prior information can never degrade CS recoverability, but may significantly enhance it.
- 1st-order methods can be fast thanks to solution sparsity (finite convergence, rates depending on sub-Hessians).
- TV models can be solved quickly if structures exploited.
- Continuation is necessary to make algorithms practical.
- Rice has a long tradition in optim. algorithms/software.



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#### Software FPC and FTVd available at:

http://www.caam.rice.edu/~optimization/L1

## **Thank You!**



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