

## Sample problems for the final exam: Some solutions

Any problem may be altered or replaced by a different one!

**Problem 1** The number 63000 has how many positive divisors?

**Solution:** 96.

First we decompose the given number into a product of primes:

$$63000 = 63 \cdot 10^3 = (7 \cdot 9) \cdot (2 \cdot 5)^3 = 2^3 \cdot 3^2 \cdot 5^3 \cdot 7.$$

An integer  $n \geq 2$  is a divisor of 63000 if and only if its prime factorisation is part of the above prime factorisation, that is, if  $n = 2^{m_1} 3^{m_2} 5^{m_3} 7^{m_4}$ , where  $0 \leq m_1 \leq 3$ ,  $0 \leq m_2 \leq 2$ ,  $0 \leq m_3 \leq 3$ , and  $0 \leq m_4 \leq 1$ . Note that the divisor  $n = 1$  admits this representation as well, with  $m_1 = m_2 = m_3 = m_4 = 0$ . By the Unique Factorisation Theorem, the quadruple  $(m_1, m_2, m_3, m_4)$  is uniquely determined by  $n$ . Thus we have a one-to-one correspondence between positive divisors of 63000 and elements of a Cartesian product  $\{0, 1, 2, 3\} \times \{0, 1, 2\} \times \{0, 1, 2, 3\} \times \{0, 1\}$ . The Cartesian product has  $4 \cdot 3 \cdot 4 \cdot 2 = 96$  elements.

**Problem 2** Solve a system of congruences (find all solutions):

$$\begin{cases} x \equiv 2 \pmod{5}, \\ x \equiv 3 \pmod{6}, \\ x \equiv 6 \pmod{7}. \end{cases}$$

**Solution:**  $x = 27 + 210k$ ,  $k \in \mathbb{Z}$ .

The moduli 5, 6 and 7 are pairwise coprime. By the generalized Chinese Remainder Theorem, all solutions of the system form a single congruence class modulo  $5 \cdot 6 \cdot 7 = 210$ . It remains to find a particular solution. One way to do this is to represent 1 as an integral linear combination of  $6 \cdot 7 = 42$ ,  $5 \cdot 7 = 35$  and  $5 \cdot 6 = 30$  (note that 1 is the greatest common divisor of these numbers). Let us apply the generalized Euclidean algorithm (in matrix form) to 42, 35 and 30:

$$\left( \begin{array}{ccc|c} 1 & 0 & 0 & 42 \\ 0 & 1 & 0 & 35 \\ 0 & 0 & 1 & 30 \end{array} \right) \rightarrow \left( \begin{array}{ccc|c} 1 & 0 & -1 & 12 \\ 0 & 1 & 0 & 35 \\ 0 & 0 & 1 & 30 \end{array} \right) \rightarrow \left( \begin{array}{ccc|c} 1 & 0 & -1 & 12 \\ -2 & 1 & 2 & 11 \\ 0 & 0 & 1 & 30 \end{array} \right) \rightarrow \left( \begin{array}{ccc|c} 3 & -1 & -3 & 1 \\ -2 & 1 & 2 & 11 \\ 0 & 0 & 1 & 30 \end{array} \right).$$

From the first row of the last matrix we read off that  $3(6 \cdot 7) - 1(5 \cdot 7) - 3(5 \cdot 6) = 1$ . Then one of the solutions is  $x = 2(3 \cdot 6 \cdot 7) + 3(-1 \cdot 5 \cdot 7) + 6(-3 \cdot 5 \cdot 6) = 252 - 105 - 540 = -393$ . Another solution is  $-393 + 2 \cdot 210 = 27$ .

**Problem 3** Find all integer solutions of a system

$$\begin{cases} 2x + 5y - z = 1, \\ x - 2y + 3z = 2. \end{cases}$$

**Solution:**  $x = -3 - 13k$ ,  $y = 2 + 7k$ ,  $z = 3 + 9k$ , where  $k \in \mathbb{Z}$ .

First we solve the second equation for  $x$  and substitute it into the first equation:

$$\begin{cases} 2x + 5y - z = 1, \\ x - 2y + 3z = 2 \end{cases} \iff \begin{cases} 2(2y - 3z + 2) + 5y - z = 1, \\ x = 2y - 3z + 2 \end{cases} \iff \begin{cases} 9y - 7z = -3, \\ x = 2y - 3z + 2. \end{cases}$$

For any integer solution of the equation  $9y - 7z = -3$ , the number  $y$  is a solution of the linear congruence  $9y \equiv -3 \pmod{7}$ . Solving the congruence, we obtain

$$9y \equiv -3 \pmod{7} \iff 2y \equiv 4 \pmod{7} \iff y \equiv 2 \pmod{7}.$$

Hence  $y = 2 + 7k$ , where  $k \in \mathbb{Z}$ . Now we find  $z$  and  $x$  by back substitution:  $z = (9y + 3)/7 = (9(2 + 7k) + 3)/7 = 3 + 9k$  and  $x = 2y - 3z + 2 = 2(2 + 7k) - 3(3 + 9k) + 2 = -3 - 13k$ . Note that  $z$  and  $x$  are integers for all  $k \in \mathbb{Z}$ .

**Problem 4** You receive a message that was encrypted using the RSA system with public key  $(55, 27)$ , where 55 is the base and 27 is the exponent. The encrypted message, in two blocks, is  $4/7$ . Find the private key and decrypt the message.

**Solution:** The private key is  $(55, 3)$ , the decrypted message is  $9/13$ .

First we find  $\phi(55)$ . The prime factorisation of 55 is  $5 \cdot 11$ , hence

$$\phi(55) = \phi(5)\phi(11) = (5 - 1)(11 - 1) = 40.$$

The private key is  $(55, \beta)$ , where the exponent  $\beta$  is the inverse of 27 (the exponent from the public key) modulo  $\phi(55) = 40$ . It is easy to find by inspection that  $\beta = 3$  (as  $3 \cdot 27 = 81 \equiv 1 \pmod{40}$ ). The standard way to find  $\beta$  is to apply the Euclidean algorithm (in matrix form) to 27 and 40:

$$\left( \begin{array}{cc|c} 1 & 0 & 27 \\ 0 & 1 & 40 \end{array} \right) \rightarrow \left( \begin{array}{cc|c} 1 & 0 & 27 \\ -1 & 1 & 13 \end{array} \right) \rightarrow \left( \begin{array}{cc|c} 3 & -2 & 1 \\ -1 & 1 & 13 \end{array} \right).$$

From the first row we read off that  $3 \cdot 27 - 2 \cdot 40 = 1$ , which implies that 3 is the inverse of 27 modulo 40.

Now that we know the private key, the decrypted message is  $b_1/b_2$ , where  $b_1 \equiv 4^3 \pmod{55}$ ,  $b_2 \equiv 7^3 \pmod{55}$ , and  $0 \leq b_1, b_2 < 55$ . We find that  $b_1 = 9$ ,  $b_2 = 13$ .

**Problem 5** Consider a relation  $\sim$  on the symmetric group  $S(n)$  defined as follows. For any  $\pi, \sigma \in S(n)$  we let  $\pi \sim \sigma$  if and only if  $\pi$  is conjugate to  $\sigma$ , which means that  $\pi = \tau\sigma\tau^{-1}$  for some permutation  $\tau \in S(n)$ . Show that  $\sim$  is an equivalence relation.

We have to show that the relation  $\sim$  is reflexive, symmetric, and transitive.

**Reflexivity.**  $\pi \sim \pi$  for all  $\pi \in S(n)$  since  $\pi = \tau\pi\tau^{-1}$  holds for  $\tau = \text{id}$  (as well as for  $\tau = \pi$ ).

**Symmetry.** Assume  $\pi \sim \sigma$ , that is,  $\pi = \tau\sigma\tau^{-1}$  for some  $\tau \in S(n)$ . Then

$$\sigma = \tau^{-1}\pi\tau = \tau^{-1}\pi(\tau^{-1})^{-1} = \tau_0\pi\tau_0^{-1},$$

where  $\tau_0 = \tau^{-1} \in S(n)$ . Hence  $\sigma \sim \pi$ .

**Transitivity.** Assume  $\pi \sim \sigma$  and  $\sigma \sim \rho$ , that is,  $\pi = \tau_1 \sigma \tau_1^{-1}$  and  $\sigma = \tau_2 \rho \tau_2^{-1}$  for some  $\tau_1, \tau_2 \in S(n)$ . Then

$$\pi = \tau_1(\tau_2 \rho \tau_2^{-1})\tau_1^{-1} = (\tau_1 \tau_2)\rho(\tau_2^{-1}\tau_1^{-1}) = (\tau_1 \tau_2)\rho(\tau_1 \tau_2)^{-1} = \tau \rho \tau^{-1},$$

where  $\tau = \tau_1 \tau_2 \in S(n)$ . Hence  $\pi \sim \rho$ .

**Problem 6** Let  $\pi = (1\ 2)(2\ 3)(3\ 4)(4\ 5)(5\ 6)$ ,  $\sigma = (1\ 2\ 3)(2\ 3\ 4)(3\ 4\ 5)(4\ 5\ 6)$ . Find the order and the sign of the following permutations:  $\pi$ ,  $\sigma$ ,  $\pi\sigma$ , and  $\sigma\pi$ .

**Solution:**  $\pi$  has order 6,  $\sigma$  has order 2,  $\pi\sigma$  and  $\sigma\pi$  have order 4. The sign of  $\sigma$  is  $+1$ , the sign of  $\pi$ ,  $\pi\sigma$  and  $\sigma\pi$  is  $-1$ .

Any transposition is an odd permutation, its sign is  $-1$ . Any cycle of length 3 is an even permutation, its sign is  $+1$ . Since the sign is a multiplicative function, we obtain that  $\text{sgn}(\pi) = (-1)^5 = -1$ ,  $\text{sgn}(\sigma) = 1^4 = 1$ , and  $\text{sgn}(\pi\sigma) = \text{sgn}(\sigma\pi) = \text{sgn}(\pi)\text{sgn}(\sigma) = -1$ .

To find the order of a permutation, we need to decompose it into a product of disjoint cycles. First we decompose  $\pi$  and  $\sigma$ :  $\pi = (1\ 2\ 3\ 4\ 5\ 6)$ ,  $\sigma = (1\ 2)(5\ 6)$ . Then we use these decompositions to decompose  $\pi\sigma$  and  $\sigma\pi$ :  $\pi\sigma = (1\ 3\ 4\ 5)$ ,  $\sigma\pi = (2\ 3\ 4\ 6)$ . The order of a product of disjoint cycles equals the least common multiple of their lengths. Therefore  $o(\pi) = 6$ ,  $o(\sigma) = 2$ , and  $o(\pi\sigma) = o(\sigma\pi) = 4$ .

**Problem 7** For any positive integer  $n$  let  $n\mathbb{Z}$  denote the set of all integers divisible by  $n$ . Does the set  $3\mathbb{Z} \cup 4\mathbb{Z} \cup 7\mathbb{Z}$  form a semigroup under addition? Does it form a group? Explain.

**Solution:** The set  $3\mathbb{Z} \cup 4\mathbb{Z} \cup 7\mathbb{Z}$  is neither a semigroup nor a group.

The set  $S = 3\mathbb{Z} \cup 4\mathbb{Z} \cup 7\mathbb{Z}$  consists of all integers divisible by at least one of the numbers 3, 4 and 7. This set is not closed under the operation of addition. For example, the numbers 4 and 7 belong to  $S$  while their sum  $4 + 7 = 11$  does not. Therefore  $S$  is neither a semigroup nor a group with respect to addition.

**Problem 8** Given a group  $G$ , an element  $g \in G$  is called central if it commutes with any element of  $G$ . The set of all central elements, denoted  $C(G)$ , is called the center of  $G$ . Prove that  $C(G)$  is a normal subgroup of  $G$ .

We need to show that the set  $C(G)$  is nonempty, closed under the group operation, and closed under taking the inverse. Clearly, the identity element  $e$  of the group  $G$  commutes with all elements of  $G$ . Hence  $e \in C(G)$ . In particular,  $C(G)$  is not empty.

Assume  $g_1, g_2 \in C(G)$ . Then for any  $h \in G$  we have  $g_1 h = h g_1$  and  $g_2 h = h g_2$ . It follows that  $(g_1 g_2) h = g_1 (g_2 h) = g_1 (h g_2) = (g_1 h) g_2 = (h g_1) g_2 = h (g_1 g_2)$ . Hence  $g_1 g_2$  is central as well.

Assume  $g \in C(G)$ . Then for any  $h \in G$  we have  $gh = hg$ . It follows that  $g^{-1} h = g^{-1} h (g g^{-1}) = g^{-1} (h g) g^{-1} = g^{-1} (g h) g^{-1} = (g^{-1} g) h g^{-1} = h g^{-1}$ . Hence  $g^{-1}$  is central as well.

**Problem 9** (i) List all cyclic subgroups of the alternating group  $A(4)$ .  
(ii) List all non-cyclic subgroups of  $A(4)$ .

**Solution:** cyclic subgroups are  $\{\text{id}\}$ ,  $\{\text{id}, (12)(34)\}$ ,  $\{\text{id}, (13)(24)\}$ ,  $\{\text{id}, (14)(23)\}$ ,  $\{\text{id}, (123), (132)\}$ ,  $\{\text{id}, (124), (142)\}$ ,  $\{\text{id}, (134), (143)\}$  and  $\{\text{id}, (234), (243)\}$ ; non-cyclic subgroups are  $\{\text{id}, (12)(34), (13)(24), (14)(23)\}$  and  $A(4)$ .

**Problem 10** All Abelian groups of order 36 form how many isomorphism classes?

**Solution:** 4.

According to the classification of finite Abelian groups, any such group is isomorphic to a direct product of cyclic groups of the form  $\mathbb{Z}_{p_1^{m_1}} \times \mathbb{Z}_{p_2^{m_2}} \times \cdots \times \mathbb{Z}_{p_k^{m_k}}$ , where  $k \geq 1$ , each  $p_i$  is a prime number, and each  $m_i$  is a positive integer. Moreover, the sequence of orders  $p_1^{m_1}, p_2^{m_2}, \dots, p_k^{m_k}$  of the cyclic groups is unique up to rearranging its terms. Note that the order of the Abelian group is the same as the order of the direct product, which equals  $p_1^{m_1} p_2^{m_2} \cdots p_k^{m_k}$ .

The prime factorisation of the number 36 is  $2^2 \cdot 3^2$ . Up to rearranging the factors, there are 4 ways to decompose it as a product of prime powers:  $36 = 2^2 \cdot 3^2 = 2 \cdot 2 \cdot 3^2 = 2^2 \cdot 3 \cdot 3 = 2 \cdot 2 \cdot 3 \cdot 3$ . It follows that all Abelian groups of order 36 form 4 isomorphism classes, represented by groups  $\mathbb{Z}_4 \times \mathbb{Z}_9$ ,  $\mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_9$ ,  $\mathbb{Z}_4 \times \mathbb{Z}_3 \times \mathbb{Z}_3$  and  $\mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_3 \times \mathbb{Z}_3$ .

**Problem 11** A linear binary coding function  $f$  is defined by a generator matrix

$$G = \begin{pmatrix} 0 & \square & 0 & 1 & 1 & 0 & 1 \\ 1 & \square & 0 & 1 & 1 & 1 & 0 \\ 0 & \square & 1 & 1 & 0 & 1 & 1 \end{pmatrix}$$

with some entries missing. Fill in the missing entries so that  $f$  can detect as many errors as possible. Explain.

**Solution:**  $G = \begin{pmatrix} 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 1 \end{pmatrix}$ .

The maximal number of errors detected by a linear binary code equals  $k-1$ , where  $k$  is the minimal weight of nonzero codewords. Suppose

$$G = \begin{pmatrix} 0 & a_1 & 0 & 1 & 1 & 0 & 1 \\ 1 & a_2 & 0 & 1 & 1 & 1 & 0 \\ 0 & a_3 & 1 & 1 & 0 & 1 & 1 \end{pmatrix},$$

where  $a_1, a_2, a_3 \in \{0, 1\}$ . Codewords of  $f$  are linear combinations of rows of the matrix  $G$  (regarded as vectors in  $\mathbb{Z}_2^7$ ). In particular,  $0a_101101$  is the first row,  $1(a_1 + a_2)00011$  is the sum of the first two rows, and  $0(a_1 + a_3)10110$  is the sum of the first and the last rows. If  $(a_1, a_2, a_3) \neq (1, 0, 0)$  then at least one of those three codewords has weight 3. On the other hand, in the case  $(a_1, a_2, a_3) = (1, 0, 0)$  all seven nonzero codewords have weight 4:  $0101101$ ,  $1001110$ ,  $0011011$ ,  $1100011$ ,  $0110110$ ,  $1010101$ , and  $1111000$ . Thus the maximal possible number of detected errors is 3, achieved for a unique choice of missing entries.

**Problem 12** The polynomial  $f(x) = x^6 + 3x^5 - 5x^3 + 3x - 1$  has how many distinct complex roots?

**Solution:** 2.

The Fundamental Theorem of Algebra implies that any polynomial  $p$  of degree  $n \geq 1$  with complex coefficients can be represented as  $p(x) = c(x - \alpha_1)(x - \alpha_2) \dots (x - \alpha_n)$ , where  $c, \alpha_1, \dots, \alpha_n \in \mathbb{C}$  and  $c \neq 0$ . The numbers  $\alpha_1, \alpha_2, \dots, \alpha_n$  are roots of  $p$ , they need not be distinct. We say that  $\alpha$  is a root of multiplicity  $k \geq 1$  if  $p(x)$  is divisible by  $(x - \alpha)^k$  but not divisible by  $(x - \alpha)^{k+1}$ . An equivalent condition is that  $p(x) = (x - \alpha)^k q(x)$  for some polynomial  $q$  such that  $q(\alpha) \neq 0$ . If this is the case then

$$p'(x) = ((x - \alpha)^k)'q(x) + (x - \alpha)^k q'(x) = k(x - \alpha)^{k-1}q(x) + (x - \alpha)^k q'(x) = (x - \alpha)^{k-1}r(x),$$

where  $r(x) = kq(x) + (x - \alpha)q'(x)$  is a polynomial and  $r(\alpha) = kq(\alpha) \neq 0$ . Hence  $\alpha$  is a root of  $p'$  of multiplicity  $k - 1$  if  $k > 1$  and not a root of  $p'$  if  $k = 1$ . We have

$$p(x) = c(x - \beta_1)^{k_1}(x - \beta_2)^{k_2} \dots (x - \beta_m)^{k_m},$$

where  $\beta_1, \dots, \beta_m$  are distinct roots of  $p$  and  $k_1, \dots, k_m$  are their multiplicities. It follows from the above that

$$\gcd(p(x), p'(x)) = (x - \beta_1)^{k_1-1}(x - \beta_2)^{k_2-1} \dots (x - \beta_m)^{k_m-1}.$$

As a consequence, the number of distinct roots of the polynomial  $p$  equals  $\deg(p) - \deg(\gcd(p, p'))$ .

To find the greatest common divisor of the polynomials  $f(x) = x^6 + 3x^5 - 5x^3 + 3x - 1$  and  $f'(x) = 6x^5 + 15x^4 - 15x^2 + 3$ , we use the Euclidean algorithm. First we divide  $f$  by  $f'$ :

$$x^6 + 3x^5 - 5x^3 + 3x - 1 = (6x^5 + 15x^4 - 15x^2 + 3)\left(\frac{1}{6}x + \frac{1}{12}\right) - \frac{5}{4}x^4 - \frac{5}{2}x^3 + \frac{5}{4}x^2 + \frac{5}{2}x - \frac{5}{4}.$$

It is convenient to replace the remainder  $r(x) = -\frac{5}{4}x^4 - \frac{5}{2}x^3 + \frac{5}{4}x^2 + \frac{5}{2}x - \frac{5}{4}$  by its scalar multiple  $\tilde{r}(x) = -\frac{4}{5}r(x) = x^4 + 2x^3 - x^2 - 2x + 1$ . Next we divide  $f'$  by  $\tilde{r}$ :

$$6x^5 + 15x^4 - 15x^2 + 3 = (x^4 + 2x^3 - x^2 - 2x + 1)(6x + 3).$$

Since  $f'$  is divisible by  $\tilde{r}$ , it follows that  $\gcd(f, f') = \gcd(f', r) = \gcd(f', \tilde{r}) = \tilde{r}$ . Thus the number of distinct complex roots of the polynomial  $f$  equals  $\deg(f) - \deg(\gcd(f, f')) = 6 - 4 = 2$ .