# Counting and Finding Real Roots of Univariate Trinomials 

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## Introduction

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Given a trinomial of the form

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where $c \neq 0,0<m<n$, and $\operatorname{gcd}(m, n)=1$ and defining

$$
r_{m, n}:=\left|\frac{n}{m^{\frac{m}{n}}(n-m)^{\frac{(n-m)}{n}}}\right|
$$

there are various cases when we compare the coefficient $c$ to $r_{m, n}$.

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- The larger root of $f$ is given by

$$
x_{h i}(c)=c^{\frac{1}{(n-m)}}\left[1-\sum_{k=1}^{\infty}\left(\frac{1}{k(n-m)^{k}} \cdot \prod_{j=1}^{k-1} \frac{k m+j(n-m)-1}{j}\right) \frac{1}{c^{\frac{k n}{(n-m)}}}\right]
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However, how can you solve without knowing how many roots there are?

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By examining the sign of the coefficients, we can determine the number of roots $f$ will have
(1) $c_{1}, c_{2}>0 \Rightarrow 0$ roots
(2) $c_{1}, c_{2}<0 \Rightarrow 1$ root
(3) $c_{1}<0$ and $c_{2}>0 \Rightarrow 1$ root

## Root Counting

When $c_{1}>0$ and $c_{2}<0, f$ can have 0,1 , or 2 roots.
By evaluating the modified A-discriminant:

$$
\equiv_{A}=\left(\frac{c_{1}}{a_{3}-a_{2}}\right)^{a_{3}-a_{2}}\left(\frac{c_{2}}{-a_{3}}\right)^{-a_{3}}\left(\frac{c_{3}}{a_{2}}\right)^{a_{2}}-1
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if this value is:
(1) $>0 \Rightarrow 2$ roots
(2) $=0 \Rightarrow 1$ root
(3) $<0 \Rightarrow 0$ roots

## Baker's Theorem

This evaluation becomes complex with large coefficients, which is why we take advantage of an application of Baker's Theorem.

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## Baker's Theorem (1966)

If $\alpha_{i} \in \mathbb{Q}_{+}, b_{i} \in \mathbb{Z}$ with $\log A_{i}:=\max \left\{h\left(\alpha_{i}\right),\left|\log \left(\alpha_{i}\right)\right|, 0.16\right\}$, $B:=\max \left\{\left|b_{i}\right|\right\}$, then

$$
\begin{gathered}
\sum_{i=i}^{m} b_{i} \log \left(\alpha_{i}\right) \neq 0 \Rightarrow \\
\log \left|\sum_{i=i}^{m} b_{i} \log \left(\alpha_{i}\right)\right|>-1.4 \cdot m^{4.5} 30^{m+3}(1+\log (B)) \prod_{i=1}^{m} \log \left(A_{i}\right)
\end{gathered}
$$

## Approximating Logarithms

## Lemma

Given any $x \in \mathbb{Q}_{+}$of height $h$, and $\ell \in \mathbb{N}$ with $\ell \geq h$ we can compute $\left\lfloor\log _{2} \max \{1, \log (x)\}\right\rfloor$ and the $\ell$ most significant bits of $\log (x)$ in time $O\left(\ell \log ^{2}(\ell)\right)$.

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Evaluating the modified A-discriminant and taking log

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We can approximate using

$$
11 L_{1}-13 L_{2}+2 L_{3}
$$

where $L_{i} \approx \log \left(\alpha_{i}\right)$ up to error $<$
$\frac{1}{10}\left(-1.4 \cdot m^{4.5} 30^{m+3}(1+\log (B)) \prod_{i=1}^{m} \log \left(A_{i}\right)\right)$

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- However these series do not produce 2 correct decimal places of until at least 20,000 terms are used.
So, we need a new series to solve this case.


## Singular Series

When $|c| \approx r_{m, n}$ the following pair of series give the roots of $f$

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$$
x_{\text {sing }}^{ \pm}=\zeta \sum_{k=0}^{\infty} \frac{\gamma_{k}}{( \pm \sqrt{(n-m) r})^{k}}(c-r)^{k / 2}
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where $\zeta=\left(\frac{m}{n-m}\right)^{1 / n}$ and $\gamma_{k} \in \mathbb{Q}[m, n]$

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We know $\gamma_{k}$ has degree $k-1$ in $(m, n)$, but an explicit formula is not currently known.

## Comparing Series



Terms needed for error $<1 / 1000$ for approximation of larger root of $f(x)=1-c x^{2}+x^{13}$ for varying values of $c$ :

| c | $x_{h i}$ | $x_{\text {sing }}^{+}$ |
| :--- | :--- | :--- |
| 1.5362173 | $>20000$ | 2 |
| 1.7 | 17 | 5 |
| 2.5 | 5 | 9 |

## Thank you!

