

MATH 614

Dynamical Systems and Chaos

**Lecture 7:**

**Symbolic dynamics (continued).**

## Symbolic dynamics

Given a finite set  $\mathcal{A}$  (an alphabet), we denote by  $\Sigma_{\mathcal{A}}$  the set of all infinite words over  $\mathcal{A}$ , i.e., infinite sequences  $\mathbf{s} = (s_1 s_2 \dots)$ ,  $s_i \in \mathcal{A}$ .

For any finite word  $w$  over the alphabet  $\mathcal{A}$ , that is,  $w = s_1 s_2 \dots s_n$ ,  $s_i \in \mathcal{A}$ , we define a **cylinder**  $C(w)$  to be the set of all infinite words  $\mathbf{s} \in \Sigma_{\mathcal{A}}$  that begin with  $w$ . The topology on  $\Sigma_{\mathcal{A}}$  is defined so that open sets are unions of cylinders. Two infinite words are considered close in this topology if they have a long common beginning.

The **shift** transformation  $\sigma : \Sigma_{\mathcal{A}} \rightarrow \Sigma_{\mathcal{A}}$  is defined by  $\sigma(s_0 s_1 s_2 \dots) = (s_1 s_2 \dots)$ . This transformation is continuous. The study of the shift and related transformations is called **symbolic dynamics**.

## Properties of the shift

- The shift transformation  $\sigma : \Sigma_{\mathcal{A}} \rightarrow \Sigma_{\mathcal{A}}$  is continuous.
- An infinite word  $\mathbf{s} \in \Sigma_{\mathcal{A}}$  is a periodic point of the shift if and only if  $\mathbf{s} = www\dots$  for some finite word  $w$ .
- An infinite word  $\mathbf{s} \in \Sigma_{\mathcal{A}}$  is an eventually periodic point of the shift if and only if  $\mathbf{s} = uwww\dots$  for some finite words  $u$  and  $w$ .
- The shift  $\sigma$  has periodic points of all (prime) periods.
- Periodic points of the shift are dense in  $\Sigma_{\mathcal{A}}$ .
- The shift transformation  $\sigma : \Sigma_{\mathcal{A}} \rightarrow \Sigma_{\mathcal{A}}$  admits a dense orbit.

## Applications of symbolic dynamics

Suppose  $f : X \rightarrow X$  is a dynamical system. Given a partition of the set  $X$  into disjoint subsets  $X_\alpha$ ,  $\alpha \in \mathcal{A}$  indexed by elements of a finite set  $\mathcal{A}$ , we can define the **itinerary map**  $S : X \rightarrow \Sigma_{\mathcal{A}}$  so that  $S(x) = (s_0 s_1 s_2 \dots)$ , where  $f^n(x) \in X_{s_n}$  for all  $n \geq 0$ .

In the case  $f$  is continuous, the itinerary map is continuous if the sets  $X_\alpha$  are **clopen** (i.e., both closed and open).

Indeed, for any finite word  $w = s_0 s_1 \dots s_k$  over the alphabet  $\mathcal{A}$  the preimage of the cylinder  $C(w)$  under the itinerary map is

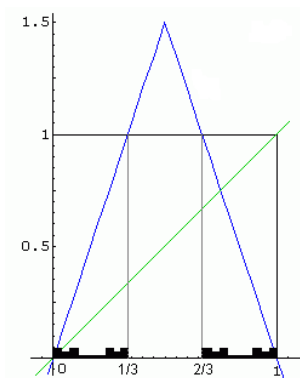
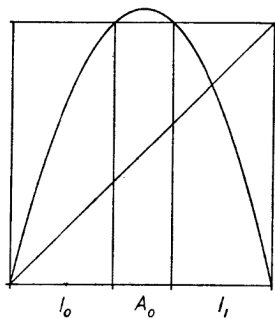
$$S^{-1}(C(w)) = X_{s_0} \cap f^{-1}(X_{s_1}) \cap \dots \cap (f^k)^{-1}(X_{s_k}).$$

## Applications of symbolic dynamics

A more general construction is to take disjoint sets  $X_\alpha$ ,  $\alpha \in \mathcal{A}$  that need not cover the entire set  $X$ . Then the itinerary map is defined on a subset of  $X$  consisting of all points whose orbits stay in the union of the sets  $X_\alpha$ .

In the case  $X$  is an interval, a partition into clopen sets is not possible. Instead, we choose the sets  $X_\alpha$  to be closed intervals with disjoint interiors. Then the itinerary map is not (uniquely) defined on a countable set.

# Examples



Any real number  $x$  is uniquely represented as  $x = k + r$ , where  $k \in \mathbb{Z}$  and  $0 \leq r < 1$ . Then  $k$  is called the **integer part** of  $x$  and  $r$  is called the **fractional part** of  $x$ . Notation:  $k = [x]$ ,  $r = \{x\}$ .

*Example.*  $f : [0, 1) \rightarrow [0, 1)$ ,  $f(x) = \{10x\}$ .

Consider a partition of the interval  $[0, 1)$  into 10 subintervals  $X_i = [\frac{i}{10}, \frac{i+1}{10})$ ,  $0 \leq i \leq 9$ . That is,  $X_0 = [0, 0.1)$ ,  $X_1 = [0.1, 0.2)$ , ...,  $X_9 = [0.9, 1)$ .

Given a point  $x \in [0, 1)$ , let  $S(x) = (s_0 s_1 s_2 \dots)$  be the itinerary of  $x$  relative to that partition. Then  $0.s_0 s_1 s_2 \dots$  is the decimal expansion of the real number  $x$ .

## Subshift

Suppose  $\Sigma'$  is a closed subset of the space  $\Sigma_{\mathcal{A}}$  invariant under the shift  $\sigma$ , i.e.,  $\sigma(\Sigma') \subset \Sigma'$ . The restriction of the shift  $\sigma$  to the set  $\Sigma'$  is called a **subshift**.

*Examples.* • Orbit closure  $\overline{O_{\sigma}^+(\mathbf{s})}$  is always shift-invariant.

- Let  $\mathcal{A} = \{0, 1\}$  and  $\Sigma'$  consists of  $(00\dots)$ ,  $(11\dots)$ , and all sequences of the form  $(0\dots 011\dots)$ . Then  $\Sigma'$  is a closed, shift-invariant set that is not an orbit closure.

- Suppose  $W$  is a collection of finite words in the alphabet  $\mathcal{A}$ . Let  $\Sigma'$  be the set of all  $\mathbf{s} \in \Sigma_{\mathcal{A}}$  that do not contain any element of  $W$  as a subword. Then  $\Sigma'$  is a closed, shift-invariant set. Any subshift can be defined this way. In the previous example,  $W = \{10\}$ .

- In the case the set  $W$  of “forbidden” words is finite, the subshift is called a **subshift of finite type**.



## Random dynamical system

Let  $f_0$  and  $f_1$  be two transformations of a set  $X$ . Consider a random dynamical system  $F : X \rightarrow X$  defined by  $F(x) = f_\xi(x)$ , where  $\xi$  is a random variable taking values 0 and 1.

The symbolic dynamics allows to redefine this dynamical system as a deterministic one. The phase space of the new system is  $X \times \Sigma_{\{0,1\}}$  and the transformation is given by

$$\mathcal{F}(x, \mathbf{s}) = (f_{s_1(\mathbf{s})}(x), \sigma(\mathbf{s})),$$

where  $s_1(\mathbf{s})$  is the first entry of the sequence  $\mathbf{s}$ .

## Substitutional dynamical systems

Given a finite set  $\mathcal{A}$ , let  $\mathcal{A}^*$  denote the set of all finite words over the alphabet  $\mathcal{A}$  (including the empty word  $\emptyset$ ).

Consider a map  $\tau : \mathcal{A} \rightarrow \mathcal{A}^* \setminus \{\emptyset\}$ , which is referred to as a **substitution rule**. This map can be extended to transformations of  $\mathcal{A}^*$  and  $\Sigma_{\mathcal{A}}$  according to the rule  $\tau(s_1s_2\dots) = \tau(s_1)\tau(s_2)\dots$  (substitutional dynamical system).

If two infinite words  $\omega, \eta \in \Sigma_{\mathcal{A}}$  have common beginning  $w$ , then the words  $\tau(\omega)$  and  $\tau(\eta)$  have common beginning  $\tau(w)$ . Note that the length of  $\tau(w)$  is not less than the length of  $w$ . It follows that the transformation  $\tau : \Sigma_{\mathcal{A}} \rightarrow \Sigma_{\mathcal{A}}$  is continuous.

## Substitutional dynamical systems

*Example.*  $\mathcal{A} = \{a, b, c, d\}$ ,

$\tau(a) = aca$ ,  $\tau(b) = d$ ,  $\tau(c) = b$ ,  $\tau(d) = c$ .

$$\tau(a) = aca,$$

$$\tau^2(a) = acabaca,$$

$$\tau^3(a) = acabacadacabaca,$$

...

Finite words  $\tau^n(a)$  “converge” to an infinite word  $\omega \in \Sigma_{\mathcal{A}}$  as  $n \rightarrow \infty$ . The infinite word  $\omega = acabacadacabaca\dots$  is a unique fixed point of the substitution  $\tau$ . This fixed point is attracting, namely,  $\tau^n(\xi) \rightarrow \omega$  for any  $\xi = a\dots \in \Sigma_{\mathcal{A}}$ .

Let  $v$  be a nonempty word over  $\{b, c, d\}$ . Then  $v\omega$  is a periodic point of period 3. For any  $\xi = va\dots \in \Sigma_{\mathcal{A}}$  the orbit  $\xi, \tau(\xi), \tau^2(\xi), \dots$  is attracted to the cycle  $v\omega, \tau(v)\omega, \tau^2(v)\omega$ .